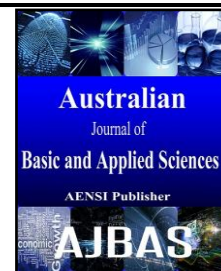




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### Impact of Genetic Algorithm Parameters on its performance for Solving Flow Shop Scheduling Problem

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#### ABSTRACT

The primary objective of flow shop scheduling is to obtain the best sequence which optimizes various objectives such as makespan, total flow time, total tardiness, or number of tardy jobs, etc. Due to the combinatorial nature of the flow shop problem (FSP) there is a lot of artificial intelligence methods proposed to solve it. The Genetic Algorithm (GA), one of these methods, is considered a valuable search algorithm capable of finding a reasonable solution in a short computational time. GA parameters, (population size, crossover probability and mutation probability) give different values that can be combined to give various GAs. In this paper we investigate the impact of the GA parameters (population size, crossover probability and mutation probability) on the quality of the GA solution in solving the flow shop scheduling problems. In this paper four population size (Ps), five crossover probability (Pc) and ten mutation probability (Pm) are investigated. The computational results show that there are significant differences among the investigated parameters on the performance of the proposed GA.

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#### INTRODUCTION

FSP is solvable to optimality in polynomial time when number of machines are limited to two,  $m = 2$ , (Johnson 1954). When the FSP enlarges as including more jobs and machines ( $m > 2$ ), it becomes a combinatorial optimization problem. It is clear that combinatorial optimization problems are NP-hard problem class, and near optimum solution techniques are preferred for such problems (Betul and Mehmet 2010).

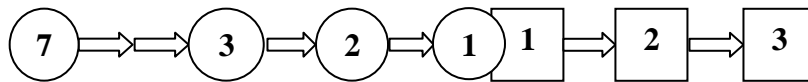
This paper is organized as follows: the next two subsections are introduction about flow shop scheduling problem, basic genetic algorithm and GA for FSP. Section 2 presents the literature review of genetic algorithm parameters (population size, crossover probability and mutation probability) in solving FSP. Section 3, present the summary of the most GA parameters used in solving FSP. Section 4 present the experimental set up, experimental design and test problem. Section 5 presents the computational results. Conclusion is presented in section 6.

#### Flow Shop Scheduling Problem (FSP):

FSP is one of the most widely studied scheduling problems in literature. It is commonly used as a benchmark for testing new heuristic algorithms, although it can hardly be applied to shop floor sequencing problems found in real production systems (Duda 2006).

In classical FSP there is a set of  $n$  jobs,  $(1, 2, \dots, n)$ , to be processed in a set of  $m$  machines,  $(1, 2, \dots, m)$ , in the same order, i.e. first in machine 1 then on machine 2 and so on, until machine  $m$ . The objective is to find a sequence for the processing of all the jobs in the machines to optimize a given criterion. In the literature, the most common criterion is the minimization of the total completion time or makespan of the schedule ( $C_{max}$ ); this is sometimes referred to as maximum flow time ( $F_{max}$ ). The processing times needed for the jobs on the machines are noted as  $p_{ij}$ . Where  $i = 1, 2, \dots, n$  and  $j = 1, 2, \dots, m$ . For instance, Figure 1 shows a sequence of 7 jobs through 3 machines.

#### Jobs Machines:



**Fig. 1:** an example of 7 jobs sequence and 3 machines.

The FSP entails a number of assumptions:

- All jobs are independent and available for processing at time 0.
- Machines are continuously available (no breakdowns).
- Each machine can only process one job at a time.
- Each job can be processed only on one machine at a time.
- Once the processing of a given job has started on a given machine, it cannot be interrupted and processing continues until completion (no preemption).
- Setup time and transportation time of the jobs are sequence independent and are included in the processing times, or ignored.
- In-process inventory is allowed. If the next machine on the sequence needed by a job is not available, the job can wait and joins the queue at that machine.
- The processing times of the jobs at the machines are known in advance.

#### **Basic Genetic Algorithm (GA):**

GA was developed by Holland in 1975 (Holland 1975). It is a search technique based on the mechanics of natural genetics and survival of the fittest. The GA object determines which individuals should survive, which should reproduce, and which should die. Since GAs are adaptive and flexible, they were shown to be successfully applied to several optimization problems (references). The basic steps of GA are as follows (Tyagi and Varshney 2012):

1. Initialize a population of binary or non-binary chromosomes.
2. Evaluate each chromosome in the population using the fitness function.
3. Select chromosome to mate (reproduction).
4. Apply genetic operators (crossover and mutation) on chromosome selected.
5. Put chromosomes produced in a temporary population.
6. If the temporary population is full, then go to step 7. Otherwise; go to step 3.
7. Replace the current population with the temporary population.
8. If termination criterion is satisfied, then quit with the best chromosome as the solution for the problem. Otherwise, go to step 2.

#### **A GA for the FSP:**

In GA method, the encoding, crossover and mutation process will give different forms. The

different forms of crossover and mutation process in GA method can be combined to give various GAs that can be used to get an optimal or near optimal solution of the FSP in reasonable CPU time. The following are generic steps for FSP GA:

*Step 1.* Based on the later review, the permutation encoding is adapted for all FSP genetic algorithms. That is, [3421 5], a chromosome, represents a job sequence where job 3 is processed first, and then job 4 is processed, and so on.

*Step 2.* In order to find the optimal solution of the problem, standard GA starts from a set of assumed or randomly generated chromosomes called initial population with size  $P_s$ , a set of solutions (chromosome) over sequence of generation.

*Step 3.* Each chromosome in the population is evaluated based on fitness criterion.

*Step 4.* Check termination criterion,  $T_c$ , if happening, stop and get the best solution. Else, reproduce a new population as follows:

*Step 5.* Two parent strings are drawn from the population according to a selection method,  $S_m$ , for reproduction. The number of copies reproduced by an individual parent is expected to be directly proportional to its fitness value.

*Step 6.* Crossover method,  $C_m$ , is used with probability  $P_c$  to recombine the two selected parents to get better offspring.

*Step 7.* Mutation method,  $M_m$ , is applied with probability  $P_m$  on the offspring generated by the  $C_m$ . It helps to preserve a reasonable level of population diversity.

*Step 8.* If the new population generated completed, go to Step#3. Else, go to Step#5.

#### **Literature Review:**

Chen *et al.* 1995 generated a GA based heuristic for FSP with makespan criterion, in which the initial population was generated by CDS (Campbell *et al.* 1970) and RA (Dannenbring's 1977). 200 problems were generated for 20 different combinations of job size and number of machines,  $n \in \{7, 10, 15, 25\}$ , and  $m \in \{4, 5, 8, 10, 15\}$ . According to generated results of trial examples, the GA default parameters are proposed as  $P_s=60, P_m=0, P_c=1$ , and 20 generation as a termination criteria. The computational results of the heuristic on a SUN workstation compared with the results of two existing heuristics, SPIRIT (Widmer and Hertz 1989) and Ho and Chang's heuristic (Ho and Chang 1991).

Reeves 1995 proposed a GA for finding the minimum makespan of  $n$ -job,  $m$ -machine permutation FSP. The initial population is

randomly generated. The default *GA* parameters used were  $Ps=30$ ,  $Pc=1$ ,  $Pm=0.8$ . On Taillard benchmarks problems (Taillard 1993), the performance of the algorithm is compared with that of a native neighborhood search technique and with a simulated annealing (*SA*) algorithm.

Murata *et al.* 1996 applied a *GA* with an objective of minimizing the makespan, and examined two hybridizations of the *GA* with other search algorithms. As test problems, they randomly generated 100 *FSP* with 20 jobs and 10 machines and 50 jobs and 10 machines. The initial population is randomly generated. Using simulations on the test problems, they found that the following specifications worked best ( $Ps=10$ ,  $Pc=1.0$ , and  $Pm=1.0$ ). Based on the default *GA* specification, the authors compared the *GA* with three search algorithms, local search (*LS*), tabu search (*TS*) and simulated annealing (*SA*).

Reeves and Yamada 1998 proposed a *GA* to minimize the makespan. The algorithm characteristics are as follow: randomly generated initial population,  $Ps=15$  and  $Pc=0.5$ . The computational are done on Taillard benchmark problems (Taillard 1993).

Tang and Liu 2002 proposed a *GA* for *FSP* with the objective to minimizing mean flow time. Two new operations are introduced into the algorithm to improve the general *GA* procedure. One replaces the worst solutions in each generation with the best solutions found in previous generation. The other improves the most promising solution through local search. Their *GA* use the following parameters,  $Ps=80$ ,  $Pc=0.99$ ,  $Pm=0.12$ , and maximum generation ( $Gmax$ ) = 1000. To evaluate the performance of the proposed *GA*, Computational experiments were carried out on a number of randomly generated problem instances,  $n \in \{50, 75, 100, 125, 150\}$  and  $m \in \{5, 10, 15, 20\}$ . Two heuristic algorithms, Ho and Chang (Ho and Chang 1991) and Rajendran (Rajendran 1993) and a basic *GA* were used for comparison.

Zdansky and Pozivil 2002 proposed a combination of genetic/tabu search algorithm (*GATS*) for *HFSP* to minimize the makespan, and compared its performance with pure *TS* and pure *GA*. The test problems are sets of randomly generated with  $n/m \in \{10/5, 12/8, 15/10\}$ . They used genetic components similar to *GA* described in (Stluka 1999). The algorithm was tested against pure *TS* and pure *GA*.

Eliter *et al.* 2004 proposed a *GA* -based heuristic for the *FSP* with makespan criterion. They used  $p_s - 1$  schedules produced by *CDS* (Campbell *et al.* 1970) method and one schedule produced by Dannenbring's method (Dannenbring 1977) to generate the initial population. Based on (Chen *et al.* 1995), the authors used  $Ps=60$ ,  $Pc=1.0$  and  $Pm=0.05$ . Based

on trial examples, they used 20 generations ( $Gmax$ ) as the termination criterion. In order to examine the effectiveness of the proposed *GA*, the performance of the algorithm is compared over 230 generated problems forming 23 different combination of jobs and machines ( $n/m = 8/5, 8/10, \dots, 35/35, 40/40$ ) with the *NEH* algorithm (Nawaz *et al.* 1983).

Iyer and Saxena 2004 proposed a *GA* for the permutation *FSP* with the objective of minimizing the makespan. They redesigned the standard *GA* implementation by using structural information from the problem. They considered five different problem dimensions,  $\frac{n}{m} \in \{\frac{10}{10}, \frac{20}{10}, \frac{49}{15}, \frac{60}{25}, \frac{100}{40}\}$ . They used two methods to simulate the matrices of processing times, one using uniform distributions and the other using normal distributions. The initial population is randomly generated. They used the following parameters based on (Bagchi and Deb 1996), ( $Ps \in \{20, 100\}$ ,  $Pc \in \{0.7, 0.9\}$  and  $Pm \in \{0.005, 0.05\}$ ). The termination of their algorithm is after 300 iterations.

Ruiz *et al.* 2006 proposed a hybrid genetic algorithm (*HGA*) that uses a simple form of local search based on the *NEH* heuristic (reference). The objective is to minimize the makespan. In the algorithm, the initial population is generated as follows: first individual based on the standard *NEH*, up to B% based on a modified *NEH*, and the remaining 100-B% randomly. The proposed experimental parameters are as follow:  $Ps \in \{20, 30, 40, 50\}$ ,  $Pc \in \{0.0, 0.1, 0.2, 0.3, 0.4\}$  and  $Pm \in \{0.0, 0.005, 0.01, 0.015\}$ . For the evaluation they used Taillard's standard benchmark (Taillard 1993) and compared the results against 11 published methods.

Sadegheih 2006 proposed a *GA*. The algorithm is tested on 8-jobs and 7-machines example for each of the combinations of the following parameter level:  $Ps \in \{20, 60, 100\}$ ,  $Pm \in \{0.001, 0.005, 0.01, 0.02\}$ , and  $Pc \in \{0.1, 0.2, 0.6, 0.8, 0.9\}$ . The author found that the proposed *GA* performance is sensitive mutation probability within a good range of values [0.005:0.01], however, outside the good range the *GA* performance deteriorates in crossover probability, the performance is insensitive to crossover rate in range [0.1:0.9] but operator is necessary. The performance is sensitive to population size in range [20:100].

Wang *et al.* 2006 proposed a Hybrid Genetic Algorithm (*HGA*) for permutation *FSP* with limited buffers with the objective to minimize the makespan. In the *HGA*, the local search based on graph model is employed. The initial population is generated as follows: one chromosome using the standard *NEH* heuristic and  $Ps -$

1 chromosomes generated randomly. They used  $P_s=50$ ,  $G_{max} = 20n * m$ . To show the effectiveness of the *HGA*, they compared the *HGA* with the simple *GA* and *TS* and *SA* with benchmarks of different scales of (Carlier 1978, Reeves 1995 and Taillard 1993).

Octavia *et al.* 2007 discussed the application of *HGA* to solve practical *FSP*. The *HGA* was run on the following sets of parameters:  $P_s=20$ ,  $(G_{max}) = 2000$ , and  $P_m=0.01$ . The initial population is generated as 1 chromosome by NEH, 1 chromosome by ascending order of job time, 1 chromosome by descending order of job time, and the rest randomly. Practical case was used as a test case to evaluate the performance of the proposed algorithm. Each test problem consisted of 120 jobs including type of product and number of unit would be produced. The proposed algorithm was compared to ant colony, genetic-tabu.

Adusumilli *et al.* 2008 proposed a *GA* for two machines *FSP* to minimize some of finishing time of arbitrary number of jobs. The default parameters are set as  $P_s=50$ ,  $(G_{max}) = 500$ ,  $P_c=0.85$ ,  $P_m=0.01$ , and the initial population is randomly. The proposed *GA* is tested on 20-jobs, 2-machines problem and compared with branch and bound technique and Johnson algorithm.

Kahraman *et al.* 2008 proposed a *GA* for *HFSP* with the objective of minimizing makespan. They gave a complete evaluation of the different parameters and operator of the algorithms using the following experiment:  $P_s=25$ ,  $P_c \in \{0.1, 0.2, 0.3, 0.4, 0.5, 0.6, 0.7, 0.8, 0.9, 1.0\}$ , and  $(P_m) \in \{0.1, 0.2, 0.3, 0.4, 0.5, 0.6, 0.7, 0.8, 0.9, 1.0\}$ . The proposed algorithm is tested on Carlier and Neron's benchmark problems (Carlier and Neron's 2000). The authors found that the best operators set are:  $(P_c) = \{0.1, 0.2, 0.3\}$ , and  $(P_m) \in \{0.1, 0.2\}$ . The results are compared with other methods that test the same problems.

Kim and Jeong 2009 proposed a flow shop scheduling with no-wait flexible lot streaming (FSS-nwFLS) using adaptive *GA* to minimize the makespan. The proposed *GA* use randomly generated initial population,  $P_s = 50$ ,  $P_c=0.5$ ,  $P_m=0.5$  and  $G_{max}=1500$ . They run 14 type of problems considering the number of products (5, 10, 15, 20, 25, 30, 50), the number of sub-lots (15, 25, 30, 45, 50, 60, 65, 75, 90, 100, 125, 150, 250). The results of the proposed *GA* are compared with other two traditional *GAs*.

Engin *et al.* 2011 proposed a *GA* based on a permutation representation of the  $n$  jobs of *HFSP* with multiprocessor task problems to minimize makespan. They proposed the following experiment: randomly generated initial population,  $P_s \in \{15, 25, 40, 50\}$ ,  $(P_c) \in \{0.1, 0.2, 0.3, 0.4, 0.5, 0.6, 0.7, 0.8, 0.9, 1.0\}$ ,  $(P_m) \in \{0.1, 0.2, 0.3, 0.4, 0.5, 0.6, 0.7, 0.8, 0.9, 1.0\}$ , stopping criterion:  $(G_{max}) = 10000$ , 25 min. CPU time). The proposed approach was tested on a set of 240 problems (Oguz 2006, Oguz and Ercan 2005 and Kahraman *et al.* 2010).

Navala 2011 proposed a review of using a *GA* based approaches in scheduling of *FMS* (flexible manufacturing system). They reviewed the most of the research papers that contributed to scheduling of *FMS* using *GAs*. They found that, in some contributions considered only one single criterion and in some multi criterion. They also observed that use of *GA* in integration with other meta heuristics like *TS*, *SA*, neural networks to determine the optimized schedule in an *FMS*.

Verma and Dhingra 2011 described multiprocessor task scheduling in the form of permutation *FSP*, which has an objective function for minimizing the makespan. They proposed the following *GA*: randomly generated initial population,  $P_s=100$ ,  $RWS$ ,  $P_c=0.6$ ,  $P_m=0.0$  and  $G_{max}=100$  as a stopping criterion. They solved an example of 15-jobs and 4-processors.

Chen *et al.* 2012 proposed a self-guided *GA* for permutation to minimize *FSP*'s makespan. In the proposed algorithm they used *TTS*,  $P_s=100$ ,  $P_c=0.6$  and  $P_m=0.3$ . They conducted extensive computational to compare the self-guided *GA* with several other algorithms using the 110 Taillard instances (Taillard 1993).

#### Summary:

Table 1 shows the summary of different *GAs* specifications used in the reviewed literatures such as population size, crossover and mutation probabilities. Also, we can notice that the following ranges:  $P_s \in [5 - 100]$ ,  $P_c \in [0 - 1.0]$ ,  $P_m \in [0 - 1.0]$ , and different problem sizes ranging from  $n \in [8 - 500]$  and  $m \in [2 - 40]$ . We used 6 selection methods, 17 crossover operators and 8 mutation operators used in designing different genetic algorithms for solving *FSP* as shown in Table 2, 3 and 4 respectively.

**Table 1:** summary of the genetic algorithms used for solving *FSP*

Author/s	Year	Criterion	$P_s$	$P_c$	$P_m$	Max. no. of Gen. ( $G_{max}$ )	Maximum size of problems ( $n, m$ )	System Specification
Chen <i>et al.</i>	1995	makespan	60	1.0	0.0	20	(25, 15)	Programmed using Fortran 77 & run on a SUN 4/490 Workstation.

Author/s	Year	Criterion	Ps	Pc	Pm	Max. no. of Gen. (Gmax)	Maximum size of problems (n, m)	System Specification
Reeves	1995	makespan	30	1.0	0.8	---	(75,20)	Programmed in Pascal, & run on a Sequent S82 Computer
Murata <i>et al.</i>	1996	makespan	5	0.5	0.5	5000	(50,10)	---
			10*	0.6	0.6			
			20	0.7	0.7			
			30	0.8	0.8			
			40	0.9	0.9			
50	1.0*	1.0*	Programmed in C, and implemented on DEC Alpha 600 5/226 Computer.					
Reeves and Yamada	1998	makespan	15	0.5	---	700	(200,20)	Pentium PC Computer
Tang and Liu	2002	mean flow time	80	0.99	0.12	1000	(150,20)	---
Zdansky and Pozivil	2002	makespan	---	---	---	100	(15,10)	Programmed in Pascal, and run on a Pentium IV(256 MB) Computer
Eliter <i>et al.</i>	2004	makespan	60	1.0	0.05	20	(40,40)	Programmed in C, and run on PCs Pentium Processors
Iyer and Saxena	2004	makespan	20	0.9*	0.05*	300	(100,40)	Programmed in (Delphi 7.0), and tested on Intel Pentium IV Processor running at 2.8 GHz with 512 MB
			100*	0.7	0.005			
Ruiz <i>et al.</i>	2006	makespan	20*	0.0	0.0	25*	(500,20)	Implemented in Turbo C++
			30	0.1	0.005	50		
			40	0.2	0.01*	75		
			50	0.3	0.015			
Sadegheih	2006	makespan	20	0.1	0.001	900	(8,7)	Coded in C++ and run on a PC with AMD Athlon 1.0 GHz CPU
			60	0.2	0.005*			
			100	0.6	0.01*			
				0.8	0.02			
				0.9				
Wang <i>et al.</i>	2006	makespan	50	---	0.2	2000	(100,20)	Coded in a C++ & developed at MIT & run under GNU g++ Compiler
Octavia <i>et al.</i>	2007	makespan	20	---	0.01	2000	n=120 factory	Implemented in Borland Delphi and run on a PC P4 Processor with 3 GHz, 512 MB
Adusumilli <i>et al.</i>	2008	makespan	50	0.85	0.01	500	(20,2)	Implemented in the Java using an IBM P 1.4 GHz Computer with 512 MB
Kahraman <i>et al.</i>	2008	makespan	25	0.1*	0.1*	72000	(15,10)	Implemented in Borland Delphi and run on a PC P4 Processor with 3 GHz, 512 MB
0.2				0.2				
0.3				0.3				
....				....				
0.9				0.9				
Kim and Jeong				1.0	1.0			

Author/s	Year	Criterion	Ps	Pc	Pm	Max. no. of Gen. (Gmax)	Maximum size of problems (n, m)	System Specification
Engin <i>et al.</i>	2009	makespan	50	0.5	0.5	1500	(50,5)	Implemented using MATLAB at command line
Verma and Dhingra	2011	makespan	15	0.1 0.2	0.1* 0.2	10000	(100,8)	Pentium 4 with 3GHz Processor and 512 MB.
			25	0.3	0.3			
Chen <i>et al.</i>	2011	makespan	40	0.9* 1.0	0.9 1.0	100	(15,4)	---
			50*	0.6	---			
Tyagi and Varshney	2012	makespan	100	0.6	0.3	300	(200,20)	
				10	0.05			
	2012	makespan	60			20	(40,20)	

\*Shows the best GA parameters (Ps, Pc and Pm).

**Table 2:** selection methods used in FSP.

#	Selection Methods (Sm)	Codes
1	Random Selection	RMS
2	Roulette Wheel Selection	RWS
3	Rank Selection	RKS
4	Tournament Selection(tour size 2)	TTS2
5	Tournament Selection(tour size 3)	TTS3
6	Tournament Selection(tour size 4)	TTS4

**Table 3:** Crossover operators methods used in FSP.

#	Crossover Operator Methods(Cm)	Codes
1	One-Point Crossover	1PX
2	Two-Point Crossover Version1	2PXV1
3	Two-point Crossover Version2	2PXV2
4	Two-Point Crossover Version3	2PXV3
5	Position Based Crossover Version1	PBXV1
6	Position Based Crossover Version2	PBXV2
7	Linear Order Crossover	LOX
8	Partially Mapped Crossover	PMX
9	Longest Common Subsequence Crossover	LCSX
10	Order Crossover	OX
11	Order Based Crossover	OBX
12	Cycle Crossover	CX
13	Similar Block Order Crossover	SBOX
14	Similar Job Order Crossover	SJOX
15	Order-Based Crossover	OBX
16	Maximal Preservative Crossover	MPX
17	Alternating Position Crossover	APX

**Table 4:** Mutation operator methods used in FSP

#	Mutation Operator Methods (Mm)	Codes
1	Adjacent Two-Job Change Mutation	AD2JM
2	Arbitrary Two-Job Change Mutation	AR2JM
3	Arbitrary Three-Job Change	AR3JM
4	Shift Mutation	SHM
5	Inversion Mutation Version1	INMV1
6	Inversion Mutation Version2	INMV2
7	Displacement Mutation	DM
8	Scramble mutation	SCM

**Table 5:** Genetic algorithm Parameters

Genetic algorithm Parameters	Values			
Crossover Probabilities ( $P_c$ )	0.3, 0.4, 0.5, 0.6, 0.7			
Mutation Probabilities ( $P_m$ )	0.001, 0.005, 0.01, 0.02, 0.05, 0.07, 0.1, 0.15, 0.2, 0.3			
Population Size ( $P_s$ )	10	20	30	50
Max. number of generation	1000	500	333	200

**Experimental Set up:**

We built a genetic algorithm based on the steps mentioned in the subsection (1.3) using *C* sharp language (Microsoft Visual Studio 2010). The six selection methods mention in Table 2, the seventeen crossover method mention in Table 3, and the eight mutation method mentioned in Table 4 are coded in the proposed *GA*.

**Experimental design:**

We design full factorial experiment as follows:

1. Based on Tables 2, 3, and 4 we have the following: 6 selection methods, 17 crossover methods and 8 mutation methods.

2. The *GA* depends also on the following other parameters: population size ( $P_s$ ), crossover probability ( $P_c$ ) and mutation probability ( $P_m$ ). We propose values for these parameters as shown in Table 5. Based on Table 5 we have the following: 4  $P_s$ , 10  $P_m$ , and 5  $P_c$ .

Therefore, the total number of combination = 6

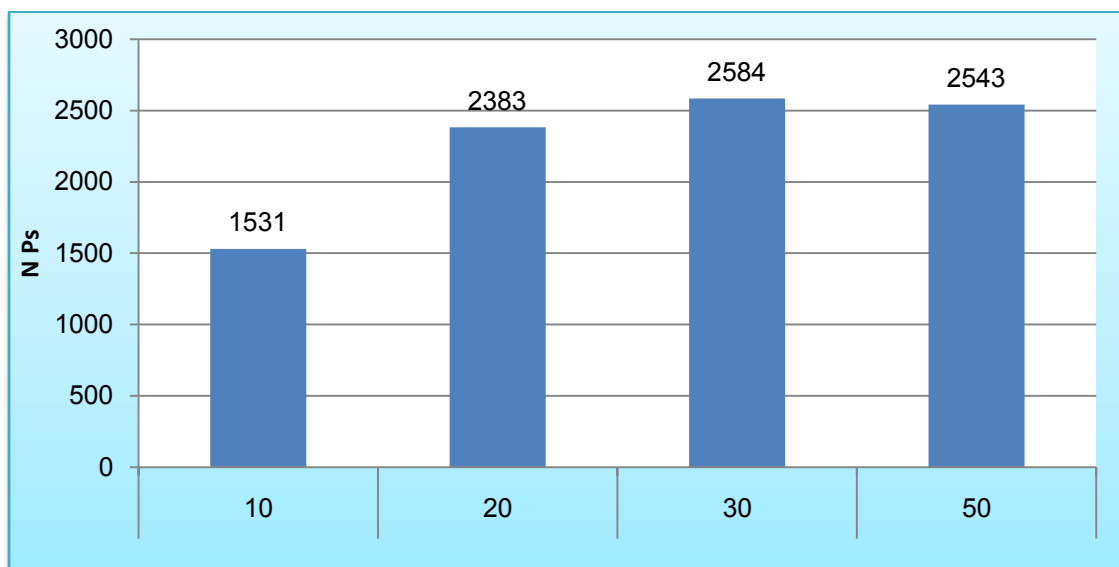
$Sm * 17 Cm * 8 Mm * 4 Ps * 5 Pc * 10 Pm = 163200$ .

**Test problem:**

We test every combination of the experiment on the first Taillard problem, ta001. The problem consists of 20 jobs and 5 machines. The processing time matrix is drawn from uniform distribution [1, 99]. The upper bound of this problem is 1278. The problem data and information can be downloaded from the OR Library [<http://people.brunel.ac.uk/~mastjib/jeb/orlib/files/flowshop1.txt>]

**Computational results:**

All computational results have been obtained on a Core 2 Duo 2.0 GHz personal computer. Each combination is run 10 times. Then, the best, average and standard deviation of makespan for the 10 are obtained.

**Fig. 2:** Population size.**Impact of the population size ( $P_s$ ):**

The contribution of the four population sizes  $P_s \in \{10, 20, 30 \text{ and } 50\}$  is determined in the figure below.

Figure 2 shows the distribution of the 9041 combinations on the population size. Applying the chi square test give (P-value equals 0.0000) which means significant differences among the population size. From the figure 2 it shows that the population size ( $P_s = 30$ ) gave best results than other population sizes, then the

(  $P_s = 50$  and 20. respectively). and the (  $P_s = 10$ ) gave worst results.

**5.2 Impact of crossover probability ( $P_c$ ):**

The contribution of the five crossover probabilities  $P_c \in \{0.3, 0.4, 0.5, 0.6 \text{ and } 0.7\}$  is determined in the figure below.

Figure 3 shows the distribution of the 9041 combinations on the crossover probability. Applying the chi square test give (P-value equals 0.0000) which means significant differences among the crossover probability. From the figure 3 it shows that

the crossover probability ( $P_c = 0.6$ ) gave best results than other crossover probabilities, then the crossover probabilities

( $P_c = 0.4, 0.7$  and  $0.5$  respectively) . And the crossover probability ( $P_c = 0.3$ ) gave the worst results.

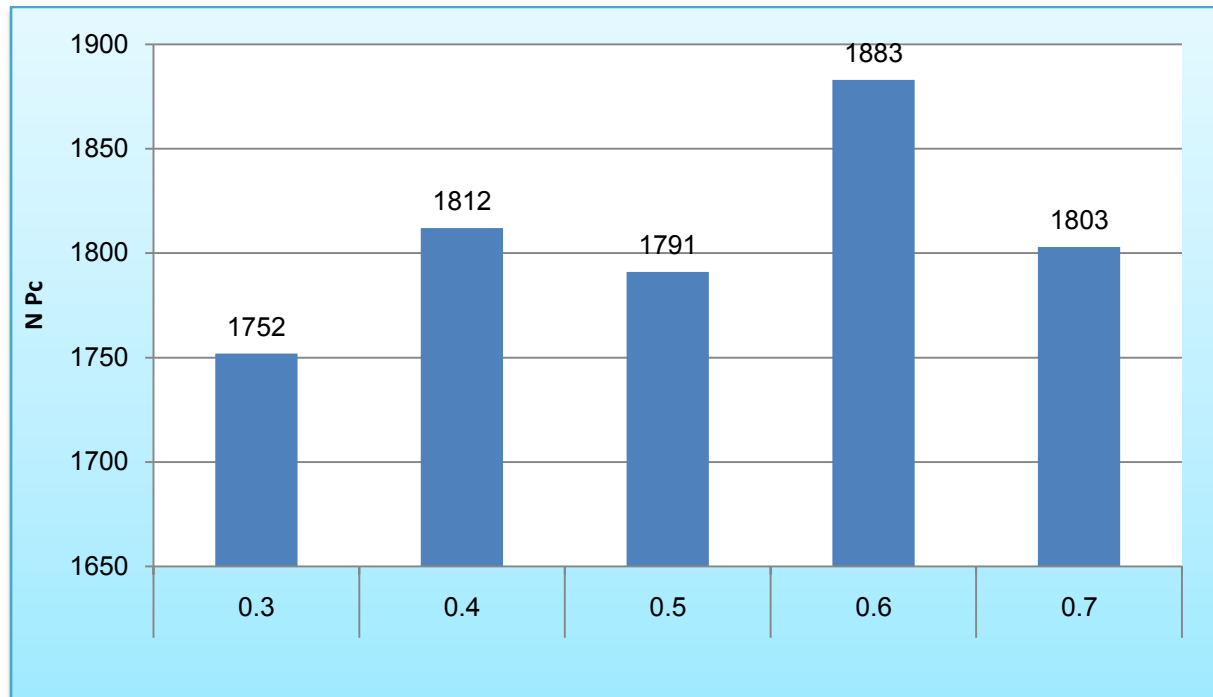


Fig. 3: Crossover probabilities.

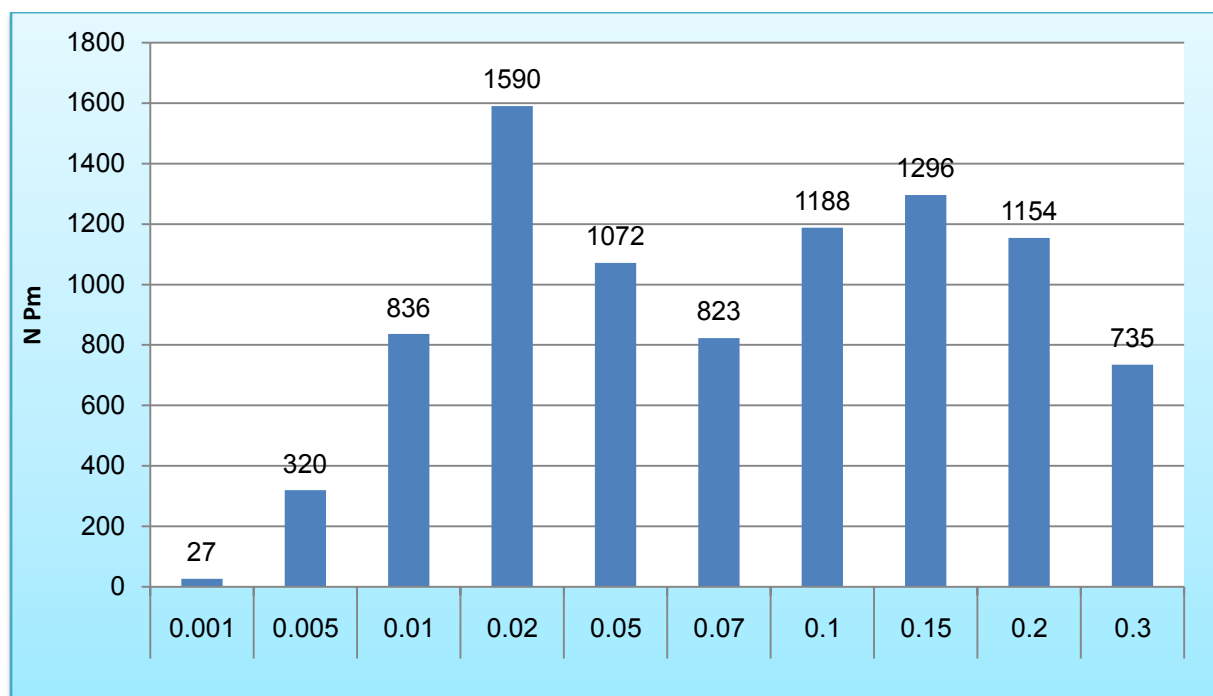


Fig. 4: Mutation probabilities.

### 5.3 Impact of the mutation probability ( $P_m$ ):

The contribution of the ten mutation probabilities

$P_m \in \{0.001, 0.005, 0.01, 0.02, 0.05, 0.07, 0.1, 0.15, 0.2 \text{ and } 0.3\}$  is determined

in the figure 4.

Figure 4 shows the distribution of the 9041 combinations on the mutation probability. Applying the chi square test give (P-value equals 0.0000) which means significant differences among the



mutation probability. From figure 4 below it shows that the mutation probability ( $P_m = 0.02$ ) gave best results than other mutation probabilities, then the mutation probabilities ( $P_m = 0.15, 0.1, 0.2$  and  $0.05$  respectively). And

the mutation probabilities gave worst results when ( $P_m = 0.001, 0.005, 0.3, 0.07,$  and  $0.01$  respectively).

**Table 6:** Default values of GA parameters.

Default Values	Type / Value
Population size ( $P_s$ )	30
Crossover Probability ( $P_c$ )	0.6
Mutation Probability ( $P_m$ )	0.02

### Conclusions:

From figures 2, 3 and 4 above, we find the default values (best GA parameters) (populationsize, crossover probability and mutation probability). (See Table 6)

It is clear that FSP problem still needs a further work such as a comparative study for all the above mentioned GA operators and parameters to see their impact on the quality of the problem solution. Moreover, the impact of problem size on the GA specification is needed.

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